

Order by Disordered Action in Swarms

Gerardo Beni
University of California Riverside

Goal

To show that
for *swarms*
lack of synchronicity
is not a hindrance
but an **advantage**

System considered

- $\sim 10^2$ - $10^{\ll 23}$ units
- \sim identical
- \sim short-range
- No centralization **No common clock**

What these systems are expected to do

- Tasks beyond the capability of a single unit
- Emergent collective 'bio-like' behavior

Collective bio-behavior

- Self-organization
- Growth
- Development
- Reproduction
- Adaptation
- Learning

ultimately depend on pattern formation

Engineering problem

- Designing systems capable of carrying out Collective bio-behavior (CBB) tasks
 - Many difficulties
 - One common difficulty in most CBB tasks:
engineering a method of pattern formation

Engineering pattern formation

- Several methods for CBB design have been proposed
- Most methods are imported from Physics and/or Computer Science

Some methods for engineering pattern formation

- Force fields
- Non-linear dynamics
- PDE
- Diffusion-reaction models
- Cellular automata
- Evolutionary algorithms
- Stochastic processes

All* methods listed
previously share one **basic
feature:**

they evolve according to a common
clock

- Explicitly
- Implicitly (as in most centralized controls)

(*) There are always some exceptions

Bio systems have **no**^(*) common clock

But, the most prevalent modes of evolution of engineered systems are:

Parallel (synchronous)

Sequential (asynchronous)

require a common clock

(*) There are always some exceptions

Swarm model based on difference equations

- Advantages
 - Quantitative
 - Predictable
- Challenge
 - Swarms are 'decentralized' and 'not synchronous'
 - Difference equations are 'centralized and synchronous'

Difference equations from differential equations

- Example: diffusion equation

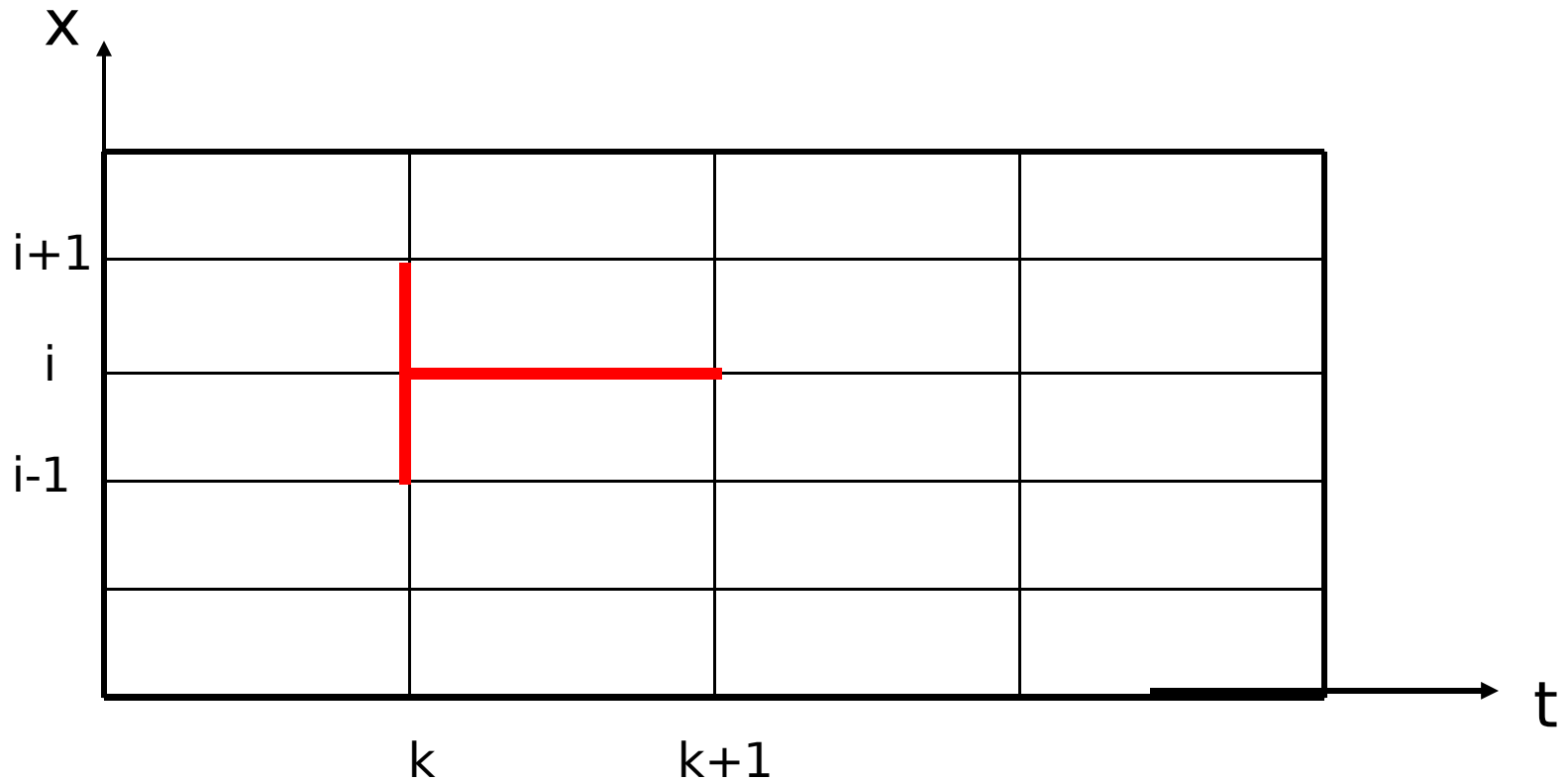
$$\frac{\partial u}{\partial t} = K \frac{\partial^2 u}{\partial x^2}$$

Diffusion difference eqn.

$$u_i^{(k+1)} - u_i^{(k)} = K\alpha (u_{i+1}^{(k)} - 2u_i^{(k)} + u_{i-1}^{(k)})$$

$$\alpha \equiv \frac{\delta t}{(\delta x)^2} = \text{Time step}/(\text{spatial step})^2$$

Solution by forward evolution



Note: Synchronicity in x common clock updates

Generalization: difference eqns. maps

- **map** is a more general concept
- A **map** is a formula that describes a new state in terms of the previous state.
- Most current research using maps focuses on **non-linear** maps (attractors, chaos, etc.)

Linear maps are also
interesting

$$\mathbf{u}^{(k+1)} = H\mathbf{u}^{(k)} + \mathbf{v}$$

\mathbf{u} , \mathbf{v} are N -dim vectors and H is a $N \times N$ matrix

Definition of Swarm as a map

- Vague definition:
 - Swarm is some sort of *self-organizing* entity
- Less vague definition:
 - Swarm is a *self-updating* map
- Math definition:
 - next slide

Swarm^(*) definition

- **A vector \mathbf{S} such that any of its component ($i=1\dots N$) updates \mathbf{S} on local time t_i according to a function s_i of \mathbf{S}**
- s_i is often more conveniently written as a function u_i of a function $v_i(\mathbf{S}, t_i) : s_i = u_i(v_i(\mathbf{S}, t_i))$
- In this way we can separate the function that singles out the components to update from the mode of updating them
- An example is the case of v_i being linear, and u_i being non linear or, e.g., stochastic.

(*) not intended to establish what a “swarm” is; this is simply an operational de

Compare **Swarms** with
a well known case of maps:

The numerical solutions of linear
systems
using iterative methods

Numerical solution of linear systems

System to solve

$$Au = b$$

Iterative solution of this system

$$u^{(k+1)} = u^{(k)} + W^{-1} (b - Au^{(k)})$$

NOTE: algorithms differ by the choice of the 'conditioning matrix' W

Classic iterative methods

- Jacobi

$$W = D$$

- Gauss-Seidel

$$W = D + L$$

- Jacobi over-relaxation JOR

$$W = \omega^{-1} D$$

- Successive over-relaxation SOR

$$W = \omega^{-1} D + L$$

Convergence

- Iterative schemes are generally recast as
$$u^{(k+1)} = Hu^{(k)} + W^{-1}b$$

$$H = I - W^{-1}A$$

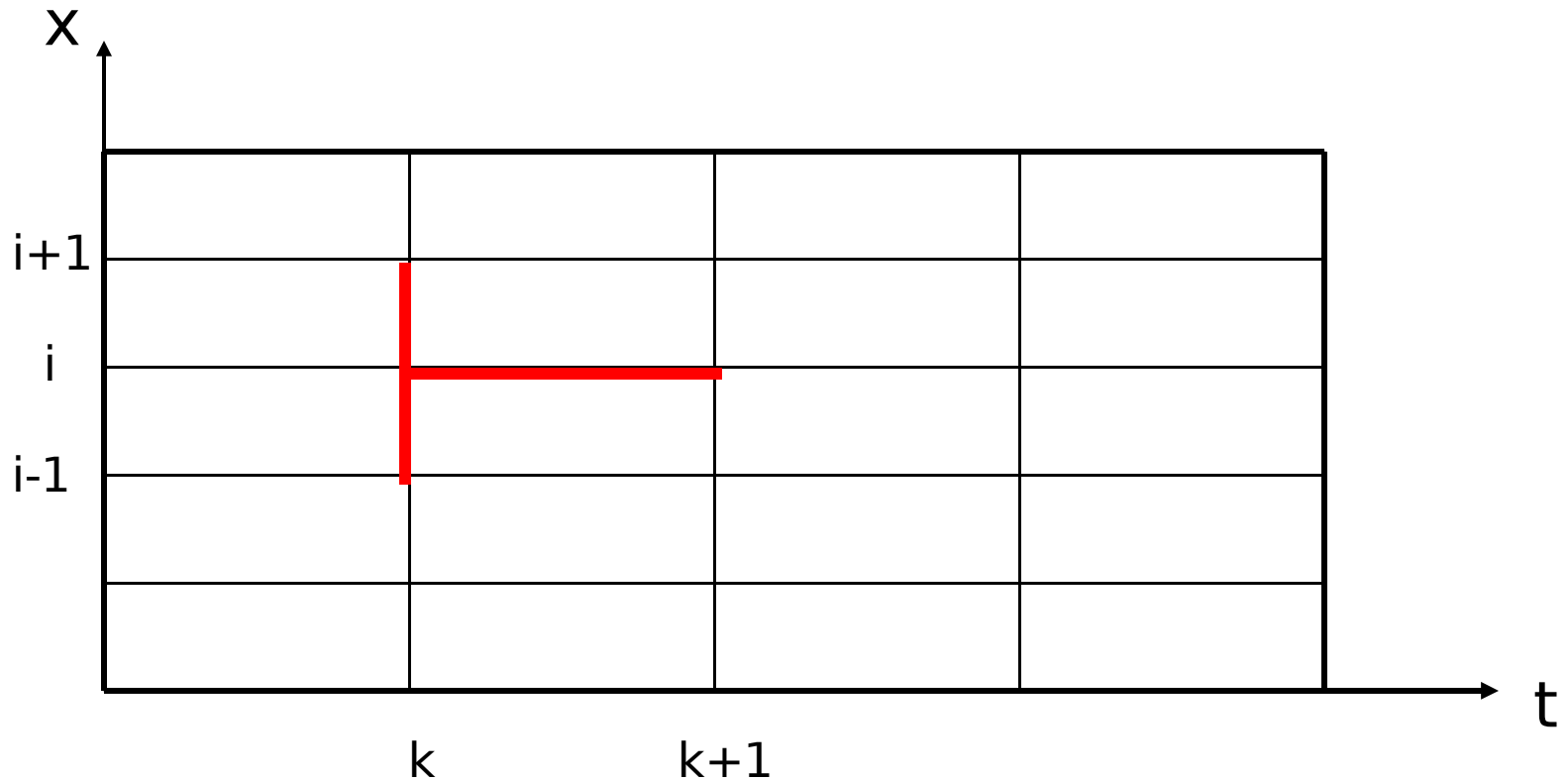
- Convergence depends only on H
- A necessary and sufficient condition is :

spectral radius of H $\rho(H) < 1$

Role of spectral radius

- Spectral radius $\rho(H)$
(max modulus of
eigenvectors)
- Convergence *iff* $\rho(H) < 1$
- Asymptotic rate of convergence $R_\infty = -\ln \rho(H)$

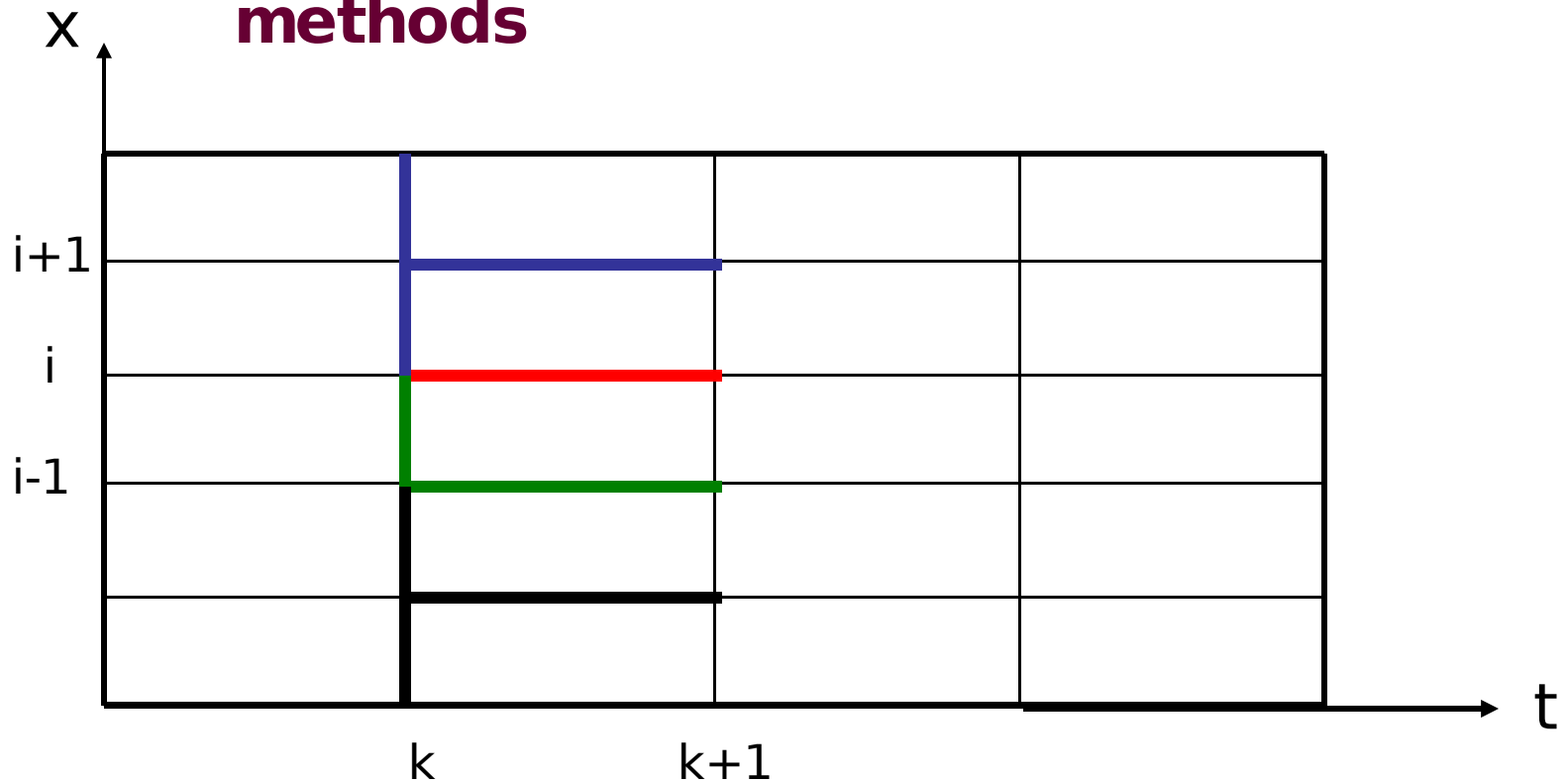
Forward evolution



Note: Synchronicity in x and 'centralized control' updates

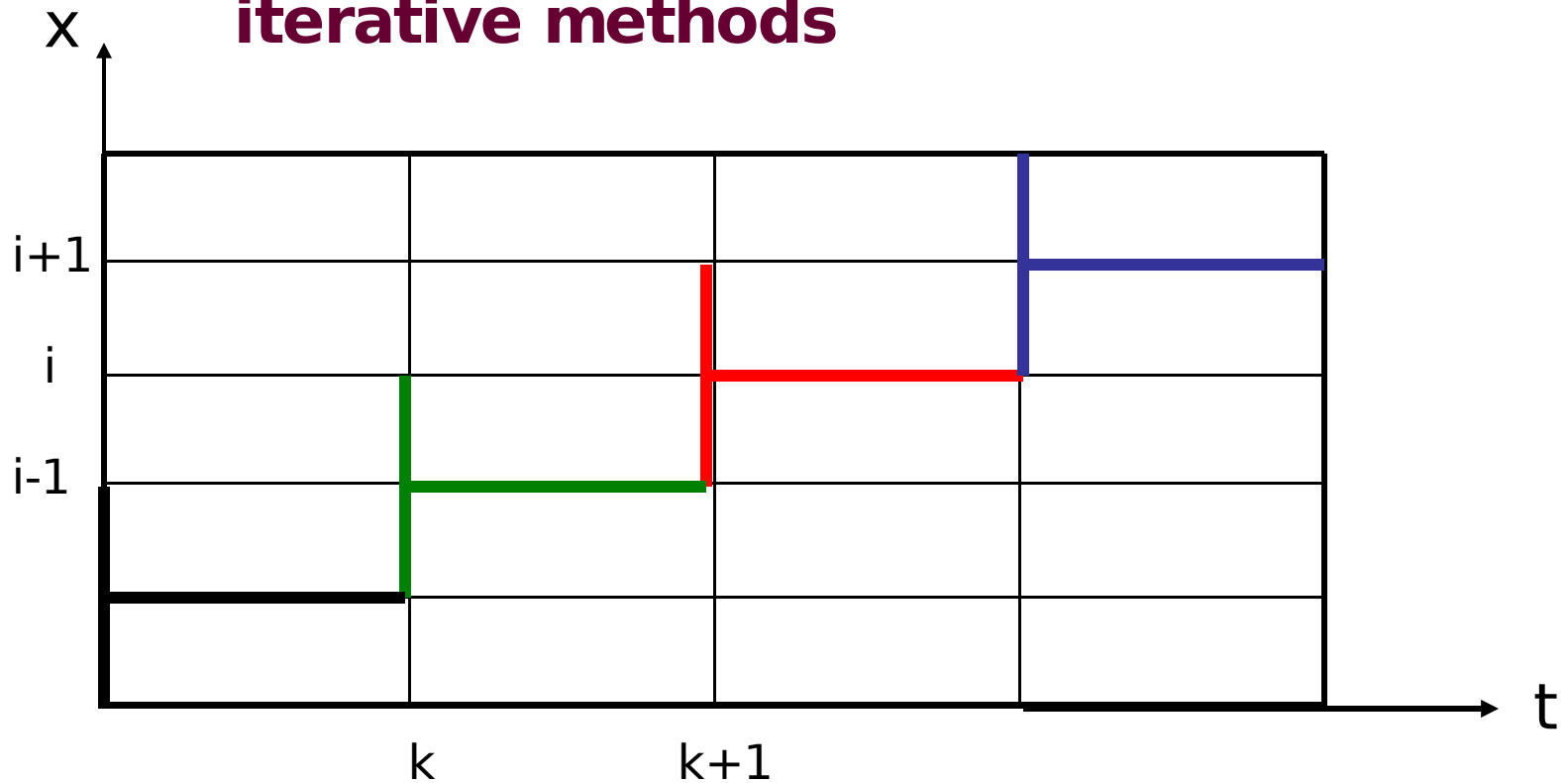
Synchronous updating

Jacobi and JOR iterative methods



Sequential updating

Gauss-Seidel and SOR iterative methods



Swarm updating methods

- Quantized over-relaxation method

$$\left| (Au - b)_i \right| < 1$$

Do nothing

$$\left| (Au - b)_i \right| \geq 1$$

$$u_i^{(k+1)} = u_i^{(k)} - 1; u_{i+1}^{(k+1)} = u_{i+1}^{(k)} + 1$$

$$\left| (Au - b)_i \right| \leq 1$$

$$u_i^{(k+1)} = u_i^{(k)} + 1; u_{i+1}^{(k+1)} = u_{i+1}^{(k)} - 1$$

- Gradient method

$$u_i^{(k+1)} = u_i^{(k)} - \epsilon \operatorname{sign} \left(\frac{\partial E}{\partial u_i^{(k)}} \right)$$

$$E = \sum_i (Au - b)_i^2$$

Quantized over-relaxation

- + converges in many cases
- - proof restricts norm of A
- - proof is only a sufficient condition

- - Only 1-dim cases (practically) solvable

Gradient Method

- +Converges
- -Needs one level more neighbors than interacting units
- -slow

Modified iterative scheme

- Iterative scheme for solving $Au = b$

$$u^{(k+1)} = Hu^{(k)} + W^{-1}b$$

$$H = I - W^{-1}A$$

- Modified iterative scheme

Sequential updating

- Basic Idea: from the synchronous iteration matrix form 'row matrices'

H	H	H
H	H	H
H	H	H

1	0	0
0	1	0
H	H	H

1	0	0
H	H	H
0	0	1

H	H	H
0	1	0
0	0	1

Repetitiveness

Same matrices appear more than once in an updating cycle

1	0	0
0	1	0
H	H	H

H	H	H
0	1	0
0	0	1

1	0	0
0	1	0
H	H	H

1	0	0
H	H	H
0	0	1

'Randomness'

Updating cycles are not identical

1	0	0
0	1	0
H	H	H

H	H	H
0	1	0
0	0	1

1	0	0
H	H	H
0	0	1

1	0	0
H	H	H
0	0	1

H	H	H
0	1	0
0	0	1

1	0	0
0	1	0
H	H	H

Partial synchronicity

Some matrices update more than one component

1	0	0
H	H	H
0	0	1

H	H	H
0	1	0
H	H	H

1	0	0
H	H	H
H	H	H

Basic Decomposition

- From synchronous to sequential
 - a special case of asynchronous order

H	H	H
H	H	H
H	H	H

1	0	0
0	1	0
H	H	H

1	0	0
H	H	H
0	0	1

H	H	H
0	1	0
0	0	1

Asynchronous set

**Shuffle the matrices and
consider all possible
permutations**

1	0	0
0	1	0
H	H	H

H	H	H
0	1	0
0	0	1

1	0	0
H	H	H
0	0	1

Notation

H = synchronous

H''' = asynchronous (even)

H'''_(odd) = asynchronous (odd)

Proof of convergence: step 1

Asynchronicity

Case	Spectral radius of H	Spectral radius of H''''	Spectral radius of H _{odd} ''''
(i)	<1	<1	<1 or ≥ 1
(ii)	<1	≥ 1	
(iii)	≥ 1	<1	<1 or ≥ 1
(iv)	≥ 1	≥ 1	

Asynchronous conditions for convergence for case (i) and (iii) are not strict

Proof of convergence: step 2a

Repetition

- Notation

$$\begin{array}{|c|c|c|} \hline 1 & 0 & 0 \\ \hline 0 & 1 & 0 \\ \hline H & H & H \\ \hline \end{array} \equiv I_{H@3}$$

- Insert **extra matrix** in H'''

$$I_{H@N} I_{H@(N-1)} \cdots I_{H@i} \cdots I_{H@2} I_{H@1}$$

$\rho(ABCD) = \rho(DABC) = \rho(CDAB) = \dots = \rho(\text{any cyclic permutation})$ →

→ $I_{H@N} I_{H@(N-1)} \cdots I_{H@i} \cdots I_{H@2} I_{H@1} = I_{H@i} H'''$

Proof of convergence: step 2b

Repetition

Notation ∞ -norm $\equiv \|A\|_{\infty} \equiv \max \sum_{j=1}^N |a_{i,j}|$

In general $\|I_{H @ i}\|_{\infty} \geq 1$ \rightarrow **Assume⁽¹⁾** $\|I_{H @ i}\|_{\infty} = 1 \rightarrow$
 $\rightarrow \|I_{H @ i} H_p''' \|_{\infty} \leq \|I_{H @ i}\|_{\infty} \|H_p''' \|_{\infty} = \|H_p''' \|_{\infty}$

In general $\rho(H''') < 1$ ~~\rightarrow~~ $\|H_p''' \|_{\infty} < 1$ **Assume⁽²⁾**

$\|I_{H @ i} H_p''' \|_{\infty} < 1 \rightarrow \rho(I_{H @ i} H_p''') < 1$

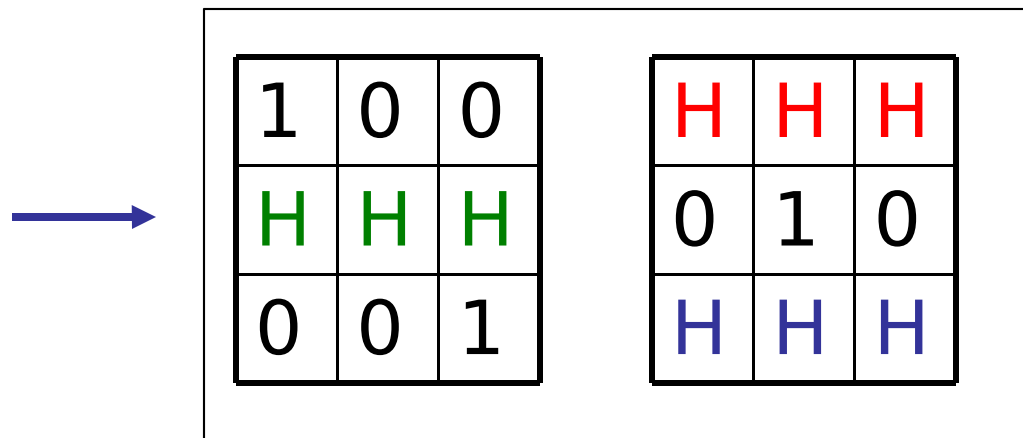
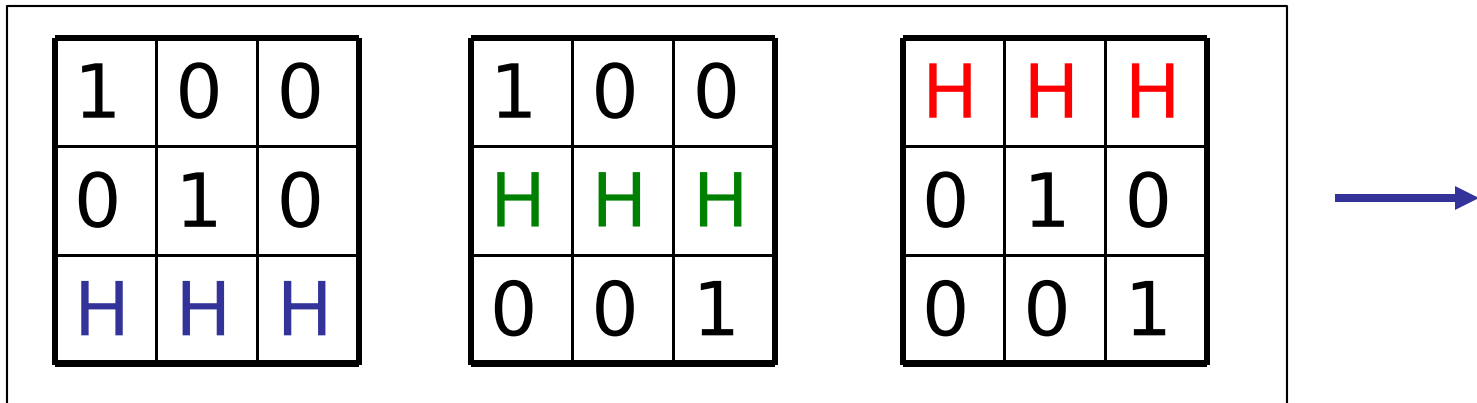
(1) And (2) not strict assumptions:

(1) works e.g. whenever H''' has zero diagonal as , e.g., in Jacobi

(2) works if ∞ -norm is replaced with L-norm (a suitable L matrix always \exists)

Proof of convergence: step 3a

Partial synchronicity



Intermediate between H and H''

Partial synchronicity

Sync/Async Transition

Case	Spectral radius of H	Spectral radius of H''''	Spectral radius of H_{odd}''''
(i)	<1	<1	<1 or ≥ 1
(ii)	<1	≥ 1	
(iii)	≥ 1	<1	<1 or ≥ 1
(iv)	≥ 1	≥ 1	

Synchronous/Async Transition

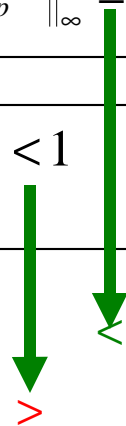
Asynchronous conditions for convergence

Repetition may lead to convergence:

In general $\|I_{H@i}\|_{\infty} \geq 1$ Assume⁽¹⁾ $\|I_{H@i}\|_{\infty} = 1 \rightarrow$

$\rightarrow \|I_{H@i}H_p'''\|_{\infty} \leq \|I_{H@i}\|_{\infty} \|H_p'''\|_{\infty} = \|H_p'''\|_{\infty}$

In general $\rho(H''') < 1$ ~~\rightarrow~~ $\|H_p'''\|_{\infty} < 1$ Assume⁽²⁾



$$\|H_p'''\|_{\infty} > 1$$

Diverges

$$\|I_{H@i}H_p'''\|_{\infty} < 1$$

Converges

Results

- **main result:** convergence under broad conditions makes pattern engineering by swarms feasible
- **Useful result:** Rate of convergence $\sim \log \rho(H)$
- **Surprising results :**
 - Sync/Async Transition
 - Convergence in cases of synchronous and/or asynchronous non convergence

Convergence

- Much broader range of conditions than previous swarm methods:
(includes, e.g., 2-dim 2nd order diff. eqs.)
- Under this broad range of conditions:
 - Partially Synchronously random-updating linear maps reach same fixed point as their synchronously updating counterparts

Rate of convergence

- Swarm $\sim r$
- SOR $\sim 2 r$
- JOR $\sim r$
- LQOR $< r$
- Swarm gradient $<< r$

$r \equiv (\text{number of cycles})^{-1}$

Clocktime vs. rate of convergence

- Average updating time of a unit : t
- Number of units: N
- Average time of an update cycle : T
- JOR (synchronous) : $T=t$
- SOR (sequential) : $T=Nt$
- Clock time to converge
- JOR : $\sim t/r$
- SOR : $\sim Nt/2r$

Synchronous/Asynchr. transition

- Swarm may **converge** even when synchronous updating does not, **if the synchronicity is reduced below a critical level.**

Reason for *convergence* in cases of synch. and/or asynchr. *non- convergence*

- Some updating orders are inaccessible in synchronous and/or asynchronous cases
- Fluctuations make accessible more updating orders, some of which lead to convergence

• Example:

$$\rho(H) = 1.304$$

$$\rho(H''_{p(\text{even})}) = 1.063$$

$$\rho(H''_{p(\text{odd})}) = 2.18$$

$$H = \begin{bmatrix} 0 & -0.45 & -0.07 \\ 0.45 & 0 & 0.07 \\ 0.07 & -0.07 & 0 \end{bmatrix}$$

Effective contraction sequences: $I_{H@3} I_{H@2} I_{H@3}$ and $I_{H@3} I_{H@1} I_{H@3}$

End

Thank you for your attention